

Enhanced Unsatisfiable Cores for QBF: Weakening Universal to Existential Quantifiers

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Intro — Some Problems with Natural Formulations as QBF

Artificial Intelligence

- Two-player games
- Variants of planning
- Many problems in knowledge representation

Formal Methods

- Verification: black box design, termination check
- Synthesis

Prototypical PSPACE-complete problem.

Introduction — Unsatisfiable Cores

Idea

- Part of an unsatisfiable formula that is by itself unsatisfiable.
- Typically obtained by syntactic weakening.

Some Applications

- Causes and explanations of unsatisfiability.
(Extends to (un)wanted implications.)
- Via duality: diagnoses and repairs.
- ... and many more ...

Fundamental concept in applied logic.

Introduction — Overview

Quantified Boolean Formulas in Prenex Conjunctive Normal Form

$$\underbrace{Q_1 p_1 \dots Q_n p_n}_{\text{prefix}} \cdot \underbrace{(l_{1,1} \vee \dots \vee l_{1,n_1}) \wedge \dots \wedge (l_{m,1} \vee \dots \vee l_{m,n_m})}_{\text{matrix: (propositional) CNF}}$$

$Q_i \in \{\exists, \forall\}$, p_i Boolean variables, $l_{i,j}$ literals over p_1, \dots, p_n .

Existing notion of unsatisfiable cores: remove clauses from matrix

$$\forall p . (p) \wedge (\neg p) \rightsquigarrow \forall p . (p) \wedge (\neg p), \quad \forall p . (p), \quad \forall p . (\neg p).$$

This paper: additionally weaken \forall to \exists

$$\forall p . (p) \wedge (\neg p) \rightsquigarrow \dots, \quad \exists p . (p) \wedge (\neg p).$$

\Rightarrow More causes/explanations of unsatisfiability. (Transfers to repairs.)

UCs for QBF in PCNF — Definitions

Let $\Pi.C$ be a QBF in PCNF.

Definition (C-, Q-, and QC-Core)

- C-Core** Remove 0 or more clauses from the matrix C [YM05].
- Q-Core** Weaken 0 or more \forall to \exists in the prefix Π .
- QC-Core** Combined c-core and q-core.

Definition (Unsatisfiable Core)

Unsatisfiable C-/Q-/QC-Core A c-/q-/qc-core that is unsatisfiable.

Definition (Minimal Unsatisfiability)

C-Minimally Unsatisfiable Unsatisfiable and no clause can be removed from the matrix C without making the result satisfiable.

Q-Minimally Unsatisfiable Unsatisfiable and no \forall can be weakened to \exists in the prefix Π without making the result satisfiable.

UCs for QBF in PCNF — Example

Consider $\Pi.C = \forall p.(p) \wedge (\neg p)$.

C-Cores: $\Pi.C$, $\forall p.(p)$, $\forall p.(\neg p)$, $\forall p.\top$

Q-Cores: $\Pi.C$,
 $\exists p.(p) \wedge (\neg p)$

QC-Cores: $\Pi.C$, $\forall p.(p)$, $\forall p.(\neg p)$, $\forall p.\top$,
 $\exists p.(p) \wedge (\neg p)$, $\exists p.(p)$, $\exists p.(\neg p)$, $\exists p.\top$

Unsatisfiable cores are red, satisfiable ones are green.

A2AECC — Q- and QC-Cores as C-Cores

Let $\Pi.C$ be a QBF in PCNF.

Definition (A2AECC)

Let $\Pi' := \Pi$, $C' := C$;

For every $\forall p_i$ in Π :

Let p'_i be fresh;

Replace $\forall p_i$ with $\forall p'_i \exists p_i$ in Π' ;

Replace C' with $(p_i \rightarrow p'_i) \wedge (p_i \rightarrow p'_i) \wedge C'$;

Return $\Pi'.C'$;

Theorem (Correctness of A2AECC)

Let \tilde{P} be a subset of the universally quantified variables in Π and let \tilde{C} be the corresponding clauses added by A2AECC. Then

$\Pi.C$ with variables in \tilde{P} weakened from \forall to \exists is satisfiable
iff

A2AECC($\Pi.C$) with clauses in \tilde{C} removed is satisfiable.

Use methods and tools for c-cores to obtain q- and qc-cores.

A2AECC — Example

Consider $\Pi.C = \forall p.(p) \wedge (\neg p)$.

$A2AECC(\Pi.C) = \forall p' \exists p.(p \rightarrow p') \wedge (p' \rightarrow p) \wedge (p) \wedge (\neg p)$.

Treat $(p \rightarrow p') \wedge (p' \rightarrow p)$ as clause group [Nad10; LS08].

QC-Core of $\Pi.C$	C-Core of $A2AECC(\Pi.C)$
$\forall p.(p) \wedge (\neg p)$	$\forall p' \exists p.(p \rightarrow p') \wedge (p' \rightarrow p) \wedge (p) \wedge (\neg p)$
$\forall p.(p)$	$\forall p' \exists p.(p \rightarrow p') \wedge (p' \rightarrow p) \wedge (p)$
$\forall p. (\neg p)$	$\forall p' \exists p.(p \rightarrow p') \wedge (p' \rightarrow p) \wedge (\neg p)$
$\forall p.\top$	$\forall p' \exists p.(p \rightarrow p') \wedge (p' \rightarrow p)$
$\exists p.(p) \wedge (\neg p)$	$\forall p' \exists p.(p) \wedge (\neg p)$
$\exists p.(p)$	$\forall p' \exists p.(p)$
$\exists p. (\neg p)$	$\forall p' \exists p.(\neg p)$
$\exists p.\top$	$\forall p' \exists p.\top$

Unsatisfiable cores are red, satisfiable ones are green.

Implementation

- Extends DepQBF 6.03 [LE17], which provides some basic infrastructure to extract c-cores, with A2AECC.
- Can be used as preprocessor or unsatisfiable c-/q-/qc-core extractor.
- Optionally performs deletion-based minimization [Mar12] with clause set refinement [BLM12].

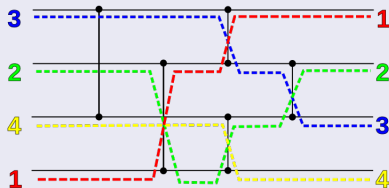
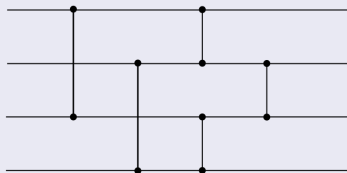
Examples

- 5342 instances from QBFLIB [GNPT]
- Interested in potential to weaken \forall to $\exists \Rightarrow$ no preprocessor

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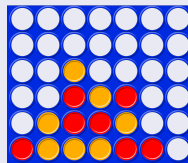
Conformant Planning: Sorting Networks [Rin07]

- Does there exist a sorting network of depth 3 for input sequences of length 6?
 - \exists plan \forall (input sequence) ...
 - Unsatisfiable core: \forall over the first number weakened to \exists .
 - No such sorting network independent of value of the first number.
 - \Rightarrow **no such sorting network of depth 3 for input sequences of length 5.**



Two-Player Games: Generalized Connect-4 [GR03]

- Can player 1 enforce a draw on a 2-by-2 board with winning lines of length 2?
 - \exists (move 1 of player 1) \forall (move 1 of player 2) ...
 - Unsatisfiable core with no \forall left.
 - Not possible, **even if player 1 had full control over the moves of player 2.**
- As before but on larger boards and with longer winning lines?
 - \exists (move 1 of player 1) \forall (move 1 of player 2) ...
 - Unsatisfiable core with a single \forall left.
 - Game is modeled [GR03] such that player 2 can play an illegal first move, thus forcing a win of player 1.
 - **Is this model of the game as intended?**



Experimental Evaluation — Overhead of UC Extraction

mode	solved instances
no unsatisfiable core	1911
unsatisfiable c-core	1830
c-minimally unsatisfiable c-core	1682
unsatisfiable q-core	1649
q-minimally unsatisfiable q-core	1139
unsatisfiable qc-core	1551
q-,c-minimally unsatisfiable qc-core	927

Related Work

- [RSMB14]: most closely related
 - introduces soft variables: may be placed at different positions in prefix, subject to preference function;
 - maximises preference function while maintaining satisfiability;
 - uses generalized version of A2AECC to reduce to weighted partial MaxSAT (we discovered our transformation independently);
 - differences:
 - makes no connection to unsatisfiable cores,
 - still satisfiable vs. still unsatisfiable,
 - always maximum vs. optionally minimal,
 - does not optimize the matrix.
- [YM05; KZ06; IJM13; LE15]: compute c-cores.
- [BLB10]: manipulates quantifiers when minimizing failure-inducing input.
- [LB11; LES16]: refer to weakening \forall to \exists as “quantifier abstraction” and “existential abstraction”.

The End

Summary

- Propose to weaken \forall to \exists in QBF unsatisfiable cores.
- Obtain additional causes of unsatisfiability.
- Implementation: enhanced UCs obtained in many instances.
- Case studies: enhanced unsatisfiable cores provide useful information.

Potential Future Work

- Understand impact of A2AECC transformation on different solvers.
- Avoid use of A2AECC transformation.
- Other logics with quantification.

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